

# Homework 5

## T1

What is the purpose of the `.END` pseudo-op? How does it differ from the `HALT` instruction?

## T2

What are the definitions of a *queue*?

## T3

The following program has an error in it. What is the error? How would you fix it?

```
.ORIG x3000
A .FILL xDEAD
B .FILL xBEEF
LD R0, A
ST R0, B
HALT
.END
```

## T4

Suppose you write two separate assembly language modules that you expect to be combined by the linker. Each module uses the label `AGAIN`, and neither module contains the pseudo-op `.EXTERNAL AGAIN`. Is there a problem using the label `AGAIN` in both modules? Why or why not?

## T5

Your friend has just written a simple program intended to calculate complements, which is as follows:

```
.ORIG x3000
; Simple program that should calculate
; complement of DATA and store the result back
LD R2, DATA
NOT R2, R2
ADD R2, R2, #1
ST R2, DATA
DATA .FILL xF001
.END
```

However, it does not seem to be reliable for some reason...

Questions:

1. What's the 2's complement of `xF001` in hex?
2. Will the program store the complement to `DATA`?
3. What will happen afterwards? Why?

Open questions (Answer if you like, but it **WILL NOT** be graded):  
What's the root cause of this phenomenon? How can we prevent this from happening?

## T6

What's the difference between pseudo-ops `.FILL`, `.BLKW` and `.STRINGZ` in LC3?

## T7

It is often useful to find the midpoint between two values. **For this problem, assume A and B are both even numbers, and A is less than B.** For example, if  $A = 2$  and  $B = 8$ , the midpoint is 5. The following program finds the midpoint of two even numbers A and B by continually incrementing the smaller number and decrementing the larger number. You can assume that `A` and `B` have been loaded with values before this program starts execution.

Your job: Insert the missing instructions.

```
.ORIG x3000
LD R0, A
LD R1, B
X ----- (a)
----- (b)
ADD R2, R2, R1
----- (c)
ADD R1, R1, #-1
----- (d)
BRnzp X
DONE ST R1,C
TRAP x25
A .BLKW 1
B .BLKW 1
C .BLKW 1
.END
```

## T8

We all know that we can achieve left-shift by adding the number to itself. For example, `ADD R0, R0, R0` will left-shift `R0` by 1 bit. However, **right-shift** is not that easy. Complete the following LC3 program so that it will right-shift `R0` by 1 bit. Note that some comments have been deleted.

```

.ORIG x3000
; Suppose R0 is already loaded with the target number
; Initialize
AND R1, R1, #0      ; Result
ADD R2, R1, #15     ; Loop var i
ADD R3, R1, #__ (a) ; 1 << (**DELETED**)
ADD R4, R1, #1      ; 1 << (15 - i)
AND R5, R5, #0      ; Temp result
; Main Loop
L  AND R5, R3, R0    ; Test bit
   BR___ (b) N      ; **DELETED**
   ADD R1, R1, R4    ; Add to result
N  ADD R3, __, __ (c) ; **DELETED**
   ADD R4, R4, R4    ; L-shift R4
   ADD __, __, __ (d) ; **DELETED**
   BRp L
; End
HALT
.END

```

## T9

The following operations are performed on a stack:

```

PUSH A
PUSH B
POP
PUSH C
POP
PUSH D
PUSH E
PUSH F
POP
PUSH G
POP
POP
POP
PUSH H

```

1. What does the stack contain after the `PUSH H` ?
2. At which point does the stack contain the most element?

Without removing the element left in the stack from the previous operations, we change this stack to a queue (the front of queue is the top of stack), and perform

```
ENQUEUE I
DEQUEUE
ENQUEUE J
ENQUEUE K
DEQUEUE
ENQUEUE L
DEQUEUE
DEQUEUE
DEQUEUE
DEQUEUE
ENQUEUE M
DEQUEUE
```

3. What does the stack contain now?

## T10

Write a function that implements another stack function, `PEEK`. `PEEK` returns the value of the top element of the stack without removing the element from the stack. The return value is stored in `R0`, so you don't need to save `R0`. `PEEK` should also do underflow error checking: if an underflow occurs, you should output the string "Stack underflow error" and halt. (Suppose the pointer of top of the stack is in `R6`, and the stack can only take up the memory space from `x3FFF` to `x3FF0`)